

Simple ambiskew polynomial rings

David Jordan

School of Mathematics and Statistics, University of Sheffield

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skew polynomial rings

base ring A

$A[x; \alpha, \delta]$

α is a ring endomorphism of A .

δ is a (left) α -derivation of A :

$$\delta(ab) = \alpha(a)\delta(b) + \delta(a)b.$$

As a group under $+$, $A[x; \alpha, \delta] = A[x]$ but

$$xr = \alpha(r)x + \delta(r).$$

Special cases: $\alpha = \text{id}_A$ ($A[x; \delta]$); $\delta = 0$ ($A[x; \alpha]$).

From now on k denotes a field.

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From now on k denotes a field.

Simple $A[x; \delta]$

A is **δ -simple** if there is no non-zero proper ideal I of A with $\delta(I) \subseteq I$.

δ is **inner** if $\exists a \in A$ such that $\delta(r) = ar - ra \forall r \in A$, **outer** otherwise.

Theorem: Let A be a k -algebra, where $\text{char } k = 0$. $A[x; \delta]$ is simple iff A is δ -simple and δ is outer.

Examples

- $A = k[y]$, $\delta = d/dy$, $xy - yx = 1$ (the first Weyl algebra).
- $A = k[y, z]$: there are examples of δ with A δ -simple, due to many people, most notably Bergman and Coutinho.

There exists a more complex simplicity criterion in characteristic p .

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There exists a more complex simplicity criterion in characteristic p .

Discussion point

What is the easiest example of a simple skew polynomial ring of the form $R[x; \alpha, \delta]$ with α outer?

There is an example due to Cozzens, reproduced over 5 pages of the 1975 book by Faith and Cozzens, of a simple ring $D[x; \alpha, \delta]$ where D is a division ring and α is injective but not surjective.

There are known simplicity results due to Lam, Leroy \pm Jain when R is a division ring and to Lam, Leroy, Leung, Matczuk when R is simple.

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ambiskew polynomial rings

The data: a k -algebra A ,

two commuting k -automorphisms of A , α and γ , $\beta := \alpha^{-1}\gamma$,
 $v \in A$ such that $va = \gamma(a)v \forall a \in A$ and $\gamma(v) = v$ (v is
 γ -normal),

$\rho \in k \setminus \{0\}$.

Form $A[y; \alpha]$ and set $\beta(y) = \rho y$ so that $\beta \in \text{Aut}_k(A[y; \alpha])$.

The **ambiskew polynomial ring** $R = R(A, \alpha, v, \rho)$ is the
iterated skew polynomial ring $A[y; \alpha][x; \beta, \delta]$, where
 $\delta(A) = 0$ and $\delta(y) = v$.

$$ya = \alpha(a)y \text{ for all } a \in A,$$

$$xa = \beta(a)x \text{ for all } a \in A \text{ and}$$

$$xy = \rho yx + v.$$

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$A = k, \alpha = \gamma = \text{id}$

- $A_1(k)$, the **first Weyl algebra**: $xy - yx = 1, \rho = \nu = 1$.
- the quantum plane: $xy = qyx, \rho = q, \nu = 0$.
- the quantised Weyl algebra $A_1^q(k)$: $xy - qyx = 1, \rho = q, \nu = 1$.

$A = k[h]: U(\mathfrak{sl}_2)$, particular examples due to Witten, Woronowicz, Podleś, Kirkman/Small, general classes due to Smith, Rueda, Le Bruyn, Benkart/Roby (down-up algebras), Cassidy/Shelton (generalized down-up algebras).

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Examples

$$A = k[K^{\pm 1}]: U_q(\mathfrak{sl}_2)$$

$$A = k[b, c]: O_q(M_2(k))$$

very recent and with combinatorial motivation

$A = k[h, K^{\pm 1}]$: augmented down-up algebras
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By iteration, higher quantized Weyl algebras.

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Casimir elements

Suppose there exists a γ -normal element $u \in A$ such that $v = u - \rho\alpha(u)$.

$z := xy - u = \rho(yx - \alpha(u))$ is such that $zy = \rho yz$, $zx = \rho^{-1}xz$, $za = \gamma(a)z$ for all $a \in A$ and $zu = uz$. Hence $zR = Rz$ and R is not simple.

If such a u exists then it is a **splitting** element, z is a **Casimir element** and R is **conformal**; otherwise R is **singular**.

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The elements $v^{(m)}$

For $m \geq 1$, $v^{(m)} := \sum_{l=0}^{m-1} \rho^l \alpha^l(v)$.

In particular $v^{(1)} = v$. $v^{(0)} := 1$.

Each $v^{(m)}$ is γ -normal.

For $m \geq 1$,

$$xy^m - \rho^m y^m x = v^{(m)} y^{m-1} \quad \text{and}$$

$$x^m y - \rho^m y x^m = x^{m-1} v^{(m)}.$$

In the conformal case, $v^{(m)} = u - \rho^m \alpha^m(u)$.

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Simplicity criterion for R

Theorem: Suppose that $\text{char } k = 0$. The ring R is simple if and only if

- 1 A is α -simple;
- 2 R is singular;
- 3 for all $m \geq 1$, $v^{(m)}$ is a unit of A .

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Examples of α -simplicity

- A a field extension of k e.g. $k = \mathbb{R}$, $A = \mathbb{C}$, $\alpha(z) = \bar{z}$.
- $A = k[t]$, $\alpha(t) = t + 1$, $\text{char } k = 0$. In this case a splitting element u always exists, of degree $1 + \deg v$ if $\rho = 1$ and $\deg v$ if $\rho \neq 1$.
- $A = k[t^{\pm 1}]$, $\alpha(t) = qt$, $q \in k^*$ not a root of unity.
- $A = kG$, $G = \langle t \rangle$ cyclic of order n , $\alpha(t) = \varepsilon t$, $\varepsilon \in k$ a primitive n th root of unity.
- A simple e.g. iterating the ambiskew construction.

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Examples of α -simplicity

- A a field extension of k e.g. $k = \mathbb{R}$, $A = \mathbb{C}$, $\alpha(z) = \bar{z}$.
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If $\text{char } k = 0$, $A_1(k)$ is, of course, simple: $v^{(m)} = m$,
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Consider \mathbb{C} as an \mathbb{R} -algebra with $\alpha(z) = \bar{z}$ and $\gamma = \text{id}_{\mathbb{C}}$. Taking $v = a + ib$ we get the \mathbb{R} -algebra R generated by i, x and y subject to the relations

$$i^2 = -1, \quad xi = -ix, \quad yi = -iy, \quad xy - \rho yx = a + ib.$$

R is singular if and only if either $\rho = 1$ and $a \neq 0$ or $\rho = -1$ and $b \neq 0$.

By the condition on $v^{(m)}$, R is simple if and only if either $\rho = 1$ and $a \neq 0$ or $\rho = -1$ and $b \neq 0$.

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Example with $A = k[t^{\pm 1}]$

If $q \in k^*$ is not a root of unity and $0 \neq n \in \mathbb{Z}$ then the k -algebra generated by $t^{\pm 1}$, x and y subject to the relations

$$\begin{aligned}yt &= qty, \\xt &= q^{-1}tx \\xy - q^{-n}yx &= t^n.\end{aligned}$$

is simple.

Example with $A = kG$, G finite cyclic

Let $\text{char } k = 0$, $n \geq 2$ and $\varepsilon \in k$ be a primitive n th root of unity.

R : k -algebra generated by t, x, y subject to

$$t^n = 1, \quad yt = \varepsilon ty, \quad xt = \varepsilon^{-1}tx, \\ xy - \rho yx = c_0 + c_1t + c_2t^2 + \dots + c_{n-1}t^{n-1}.$$

R is generically simple. For example when $n = 2$, $\varepsilon = -1$ and $\rho = 1$ the relations become

$$t^2 = 1, \quad yt = -ty, \quad xt = -tx, \\ xy - yx = c_0 + c_1t.$$

Here R is simple if and only if $c_1 \neq 0$ and $c_0 \neq \pm mc_1$ for all odd $m \geq 1$.

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Simple higher quantized Weyl algebras

Suppose $\text{char } k = 0$ and let $\Lambda = (\lambda_{i,j})$ be $n \times n$ such that $\lambda_{j,i} = \lambda_{i,j}^{-1}$ for $1 \leq i, j \leq n$. The k -algebra generated by $x_1, y_1, \dots, x_n, y_n$ subject to the relations

$$y_j y_i = \lambda_{j,i} y_i y_j, \quad 1 \leq i < j \leq n;$$

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is simple. These are iterated ambiskew polynomial rings and special cases of a more general construction.

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Removing the Casimir barrier

Suppose that R is conformal so that the Casimir element z is a barrier to simplicity. It can be removed by inverting $z = xy - u$ to form $S := R_Z := R_{\{z^i\}}$ or factoring out z to form $T := R/zR$. There are simplicity criteria for both S and T .

Simplicity criterion for S

Theorem: The ring $S = R_{\mathbb{Z}}$ is simple if and only if the following hold:

- 1 A is $\{\alpha, \gamma\}$ -simple;
- 2 there do not exist $c \in A$ and $m, j \in \mathbb{Z}$, with m and j not both 0, such that $\gamma(c) = \rho^m c, \alpha(c) = \rho^j c$ and $c\gamma^j(a) = \alpha^m(a)c \forall a \in A$;
- 3 for all $m \geq 1$, there exists n such that $u^n \in v^{(m)}A$.

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Examples

Examples with S simple:

- the quantum torus $k[x^{\pm 1}, y^{\pm 1} : xy = qyx]$, q is not a root of unity.
- A_1^q when $xy - yx$ is inverted and q is not a root of unity;
- localizations of non-simple higher quantized Weyl algebras;
- $A = k[h]$, $\alpha(h) = h + 1$, $\gamma = \text{id}$, ρ not a root of unity,
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Factor out z

Abusing notation, in $T := R/zR$,

$$\begin{aligned} ya &= \alpha(a)y \text{ for all } a \in A, \\ xa &= \beta(a)x \text{ for all } a \in A \text{ and} \\ xy &= u \quad ; \quad yx = \alpha(u). \end{aligned}$$

This is a (generalized) generalized Weyl algebra in the sense of Bavula.

Simplicity criterion for T

Theorem: T is simple if and only if

- 1 A is α -simple;
- 2 α^m is outer for all $m \geq 1$;
- 3 u is regular;
- 4 $uA + \alpha^m(u)A = A$ for all $m \geq 1$.

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